



## Trees



## What is a Tree?

Non-linear data structure

- Hierarchical arrangement of data

Has components named after natural trees

- root
- branches
- leaves

Drawn with root at the top



## Components of a Tree

**Node:** stores a data element

**Parent:** single node that directly precedes a node

- all nodes have 1 parent except root (has 0)

**Child:** one or more nodes that directly follow a node

**Ancessor:** any node which precedes a node

- itself, its parent, or an ancestor of its parent

**Descendent:** any node which follows a node

- itself, its child, or a descendent of its child



## More Tree Terminology

**Leaf (external) node:** node with no children

**Internal node:** non-leaf node

**Siblings:** nodes which share same parent

**Subtree:** a node and all its descendants

- Ignoring the node's parent, this is itself a tree

**Ordered tree:** tree with defined order of children

- enables ordered *traversal*

**Binary tree:** ordered tree with up to two children per node



## Examples of Trees

**Directory tree**

- Organizes directories and files hierarchically
- Directories are internal nodes, files are leaf nodes (usually)

**Class hierarchy**

- Object is root, other classes are descendants

**Decision tree**

- Binary tree
- Path taken determined by boolean expression

**Expression tree**

- Operators are internal nodes, variables and constants are leaf nodes



## Comparison of Tree and List

	List	Tree
<b>Start</b>	head	root
<b># before</b>	1 (prev)	1 (parent)
<b># after</b>	1 (next)	$\geq 1$ (children)



## Tree methods

`root()`: returns root  
`parent(v)`: returns parent of `v`  
`children(v)`: returns iterator of children of `v`  
`size()`: returns number of nodes  
`elements()`: returns iterator of all elements  
`positions()`: returns iterator of all positions/nodes  
`swapElements(v,w)`: swaps elements at two nodes  
`replaceElement(v,e)`: replaces element of a node

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Tree query utility methods

`isInternal(v)`: test if node is internal  
`isExternal(v)`: test if node is external  
`isRoot(v)`: test if node is root

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Depth

*Depth* of `v` is numbers of ancestors (excluding `v`)

- depth of root is 0
- depth of node is depth of parent plus 1

```

public static int depth (Tree T, Node v) {
    if (T.isRoot(v)) return 0;
    else return 1 + depth(T, T.parent(v));
} // running time?  $O(d_v)$ 
  
```

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Height

*Height* of `v` is maximum path length of subtree

- height of leaf node is 0
- height of internal node is maximum height of children + 1

—height of a tree is height of root or maximum depth of a leaf

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Naïve height algorithm

```

public static int height1(Tree T) {
    int h=0;
    PositionIterator it = T.positions();
    while (it.hasNext()) {
        Position v = it.nextPosition();
        if (T.isExternal(v))
            h = Math.max(h, depth(T,v));
    }
    return h; } // running time?  $O(n^2)$ 
  
```

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Efficient height algorithm

```

public static int height2(Tree T, Position v) {
    if (T.isExternal(v)) return 0;
    else {
        int h=0;
        PositionIterator children = T.children(v);
        while (children.hasNext())
            h = Math.max(h, height2(T, children.nextPosition()));
        return 1 + h; }
    } // running time?
  
```

$O(n)$ : each node visited once

Johns Hopkins Department of Computer Science  
 Course 600.226: Data Structures, Professor: Jonathan Cohen



## Traversal

Ordered way of *visiting* all nodes of tree  
Converts hierarchy into a linear sequence



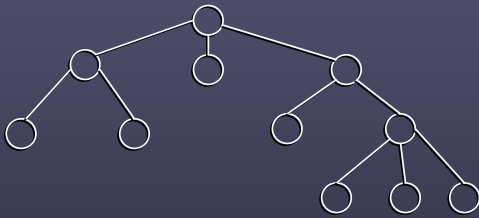
## Preorder Traversal

Visit node, then visit children

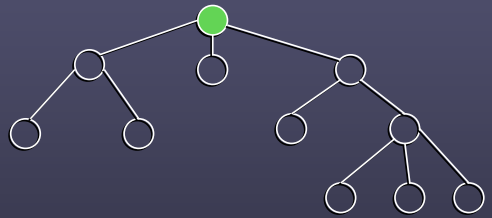
```
public void preorder(Tree T, Position v)
{
    visit(v);
    PositionIterator children = T.children(v);
    while (children.hasNext())
        preorder(children.nextPosition());
}
```



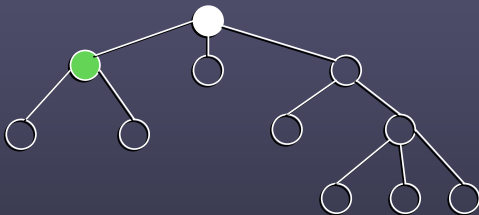
## Preorder Example



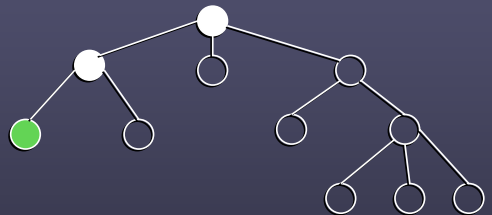
## Preorder Example



## Preorder Example

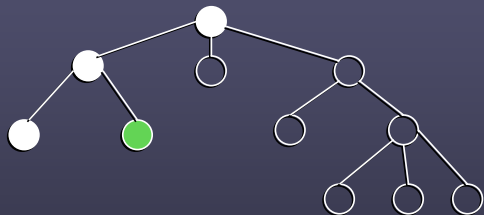


## Preorder Example





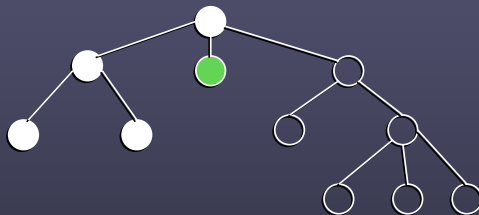
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



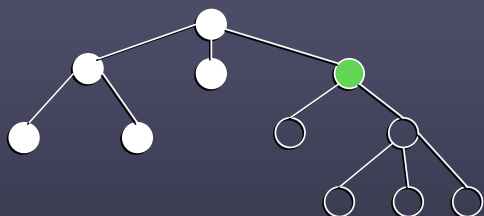
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



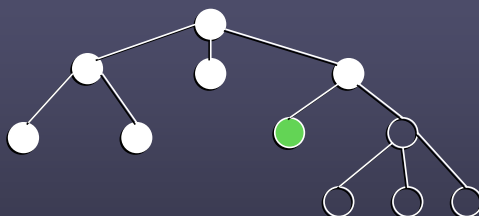
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



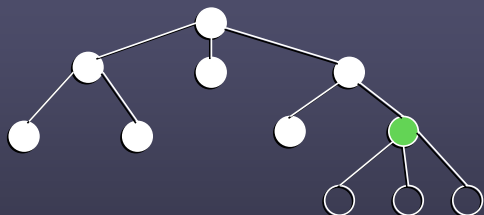
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



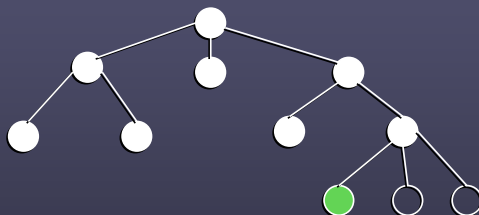
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



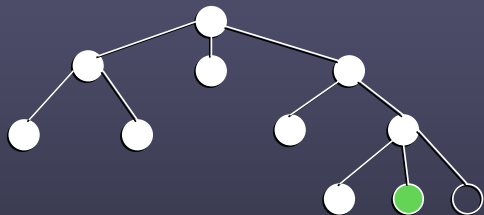
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



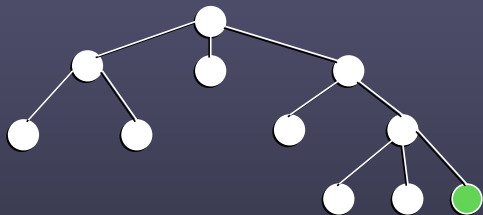
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



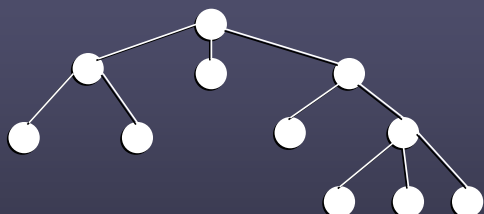
## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Preorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Postorder Traversal

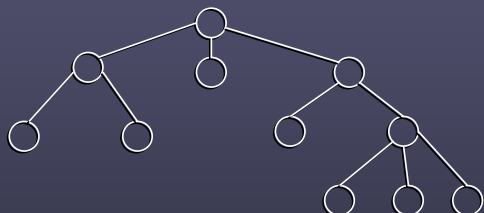
**Visit children, then visit node**

```
public void postorder(Tree T, Position v)
{
    PositionIterator children = T.children(v);
    while (children.hasNext())
        postorder(children.nextPosition());
    visit(v);
}
```

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



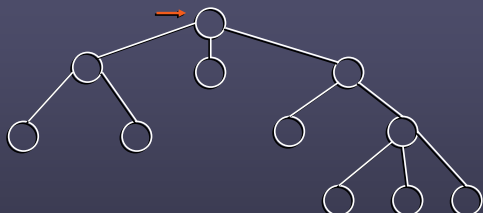
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



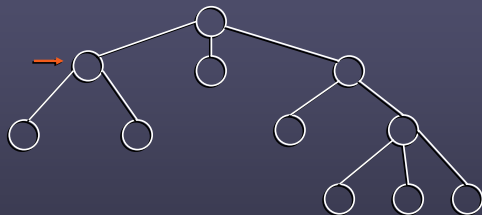
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



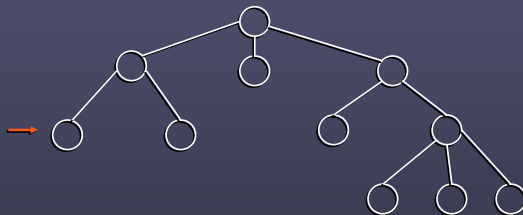
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



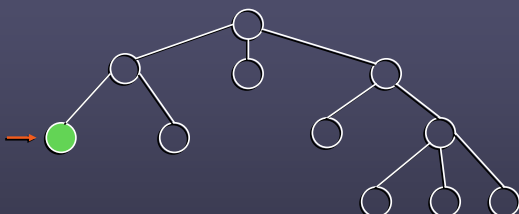
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



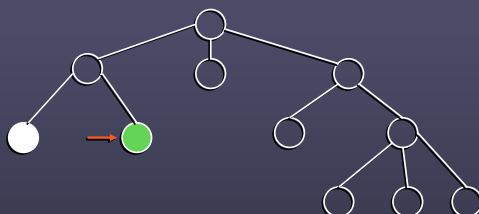
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



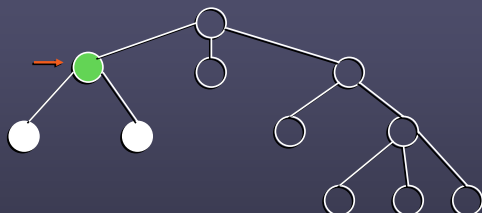
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



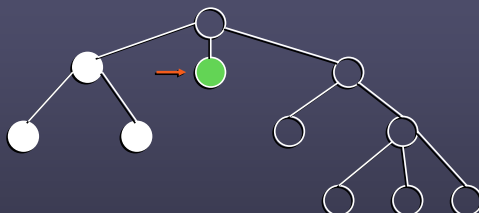
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



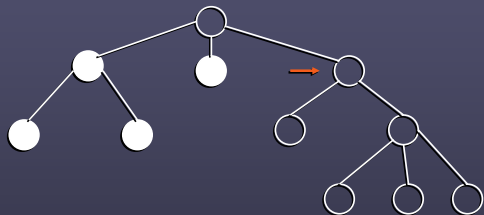
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



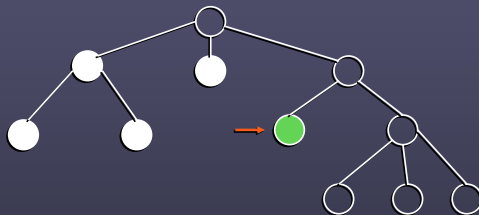
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



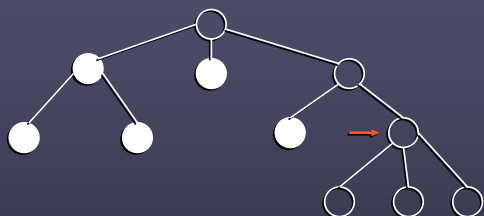
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



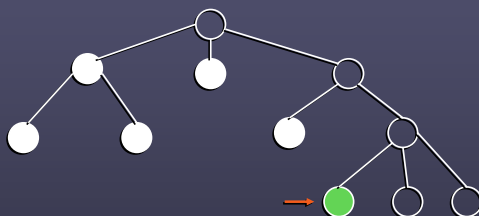
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



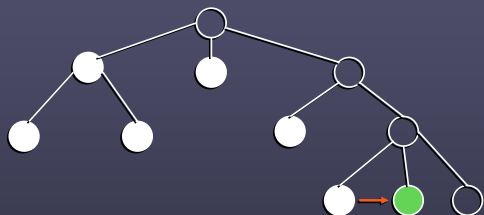
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



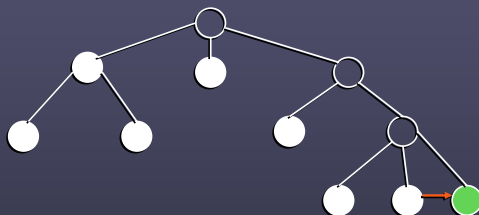
## Postorder Example



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



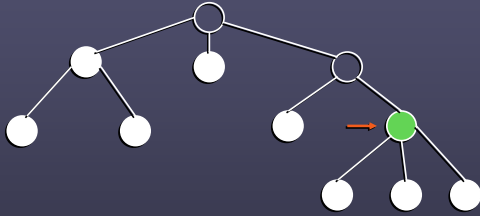
## Postorder Example



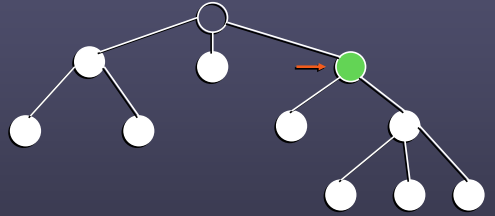
Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



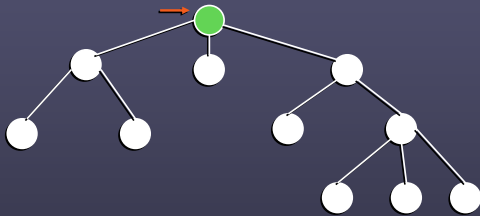
## Postorder Example



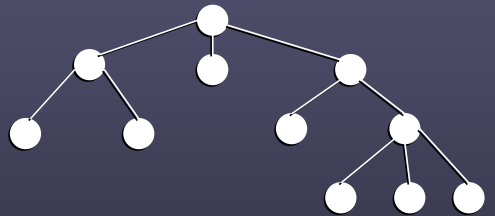
## Postorder Example



## Postorder Example



## Postorder Example

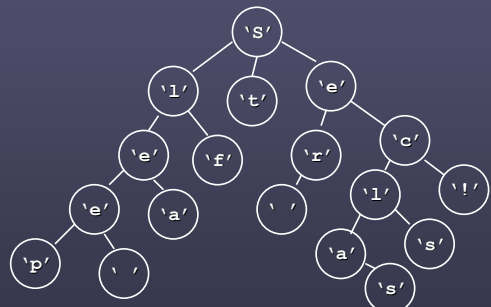


## In-class Exercises...

(Paper and pencil recommended)



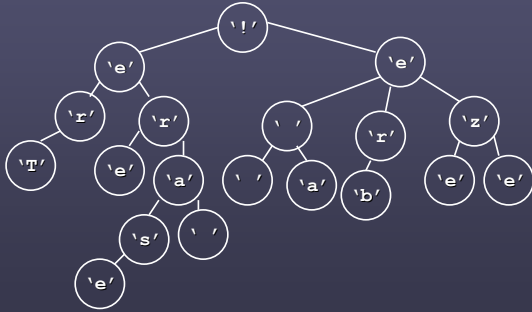
## Preorder Letter Scramble







## Postorder Letter Scramble



Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Binary Tree

Each node has no more than 2 children

- *Proper* binary tree: each node has either 0 or 2 children

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Binary Tree ADT

`leftChild(v)`: returns left child of `v`

`rightChild(v)`: returns right child of `v`

`sibling(v)`: returns sibling of `v`

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Binary Tree Traversal

Preorder: node, left, right

Postorder: left, right, node

Inorder: left, node, right

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Euler Tour Traversal

Generalizes preorder, inorder, and postorder

Visit each internal node 3 times

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Euler Tour

```
public void eulerTour(Tree T,
                      position v) {
    visitPre(T,v);
    if (T.hasLeft(v))
        eulerTour(T, T.leftChild(v));
    visitIn(T,v);
    if (T.hasRight(v))
        eulerTour(T, T.rightChild(v));
    visitPost(T,v); return; }

```

Johns Hopkins Department of Computer Science  
Course 600.226: Data Structures, Professor: Jonathan Cohen



## Template Method Pattern

Implement template or skeleton method for high-level algorithm

Extend the template's class to override lower-level methods

EulerTour example

- Override methods for visitPre, visitIn, and visitPost



## Useful Binary Tree Definitions

Level  $d$ : All nodes in a binary tree at depth  $d$

- Maximum of  $2^d$  nodes in level  $d$

*Complete* binary tree: tree of height  $h$  with  $2^h$  leaf nodes

- $2^h - 1$  internal nodes
- $2^{h+1} - 1$  total nodes



## Binary Tree Properties

(proper) Binary tree  $T$  of height  $h$

- $h+1 \leq \text{external nodes} \leq 2^h$
- $h \leq \text{internal nodes} \leq 2^h - 1$
- $2h+1 \leq \text{total nodes} \leq 2^{h+1} - 1$
- $\log(n+1) - 1 \leq h \leq (n-1)/2$
- external nodes = internal nodes + 1



## Implementing Binary Tree

Linked

- Each node references left, right, and parent as well as element

Vector-based

- Number nodes in level order
- Store nodes at rank according to number
- Storage allocated for entire complete tree